

Mike Hamburg, Rambus Cryptography Research

The STROBE protocol framework

Secure, simple, and small

What is STROBE?

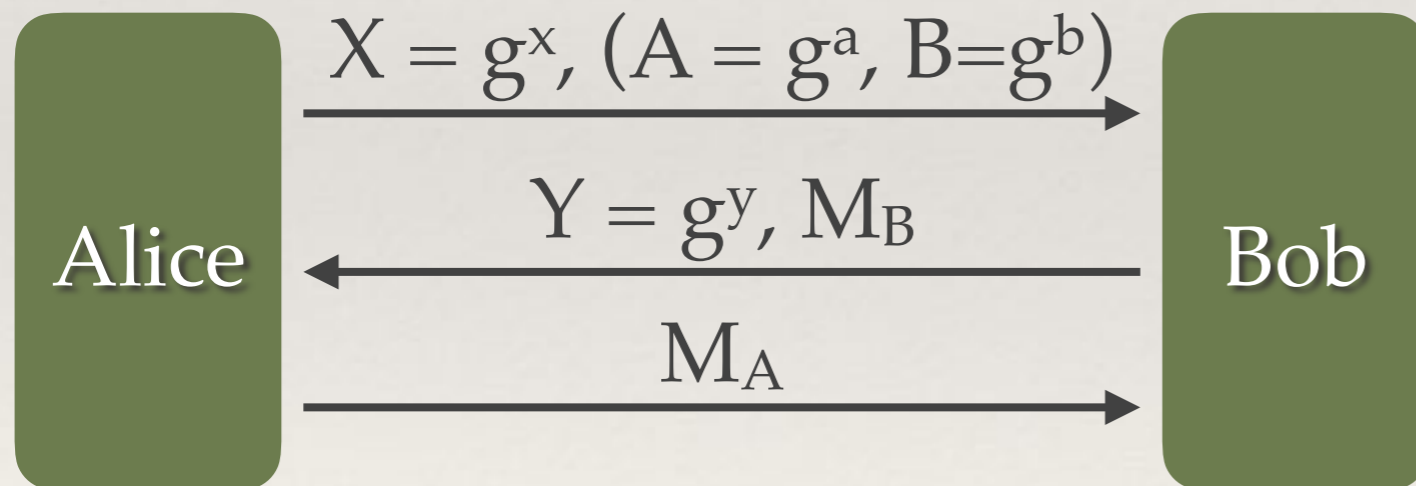
- ❖ Protocol framework with embedded focus
 - ❖ Simple protocols and handshakes
 - ❖ Encrypt, MAC, hash, sign...
- ❖ Simple, easy to analyze
- ❖ Non-terrible performance
- ❖ Can be an instance of NIST [cSHAKE]

Motivation: bespoke protocols

- ❖ Best practice: use TLS or IPSEC
- ❖ Real-world protocols have diverse requirements
 - ❖ Public key encryption / auth algorithms
 - ❖ Message flow
 - ❖ Code size and memory requirements
- ❖ Result: lots of custom protocols!
 - ❖ Design and analysis are a pain
 - ❖ Often insecure

Motivation: academic protocols

- ❖ Hash, sign, encrypt and MAC on tuples, key confirm
- ❖ [FHMQV]-C:



$$d = H(X, Y, A, B)$$

$$e = H(Y, X, A, B)$$

$$\sigma = g^{(x+da)(y+eb)}$$

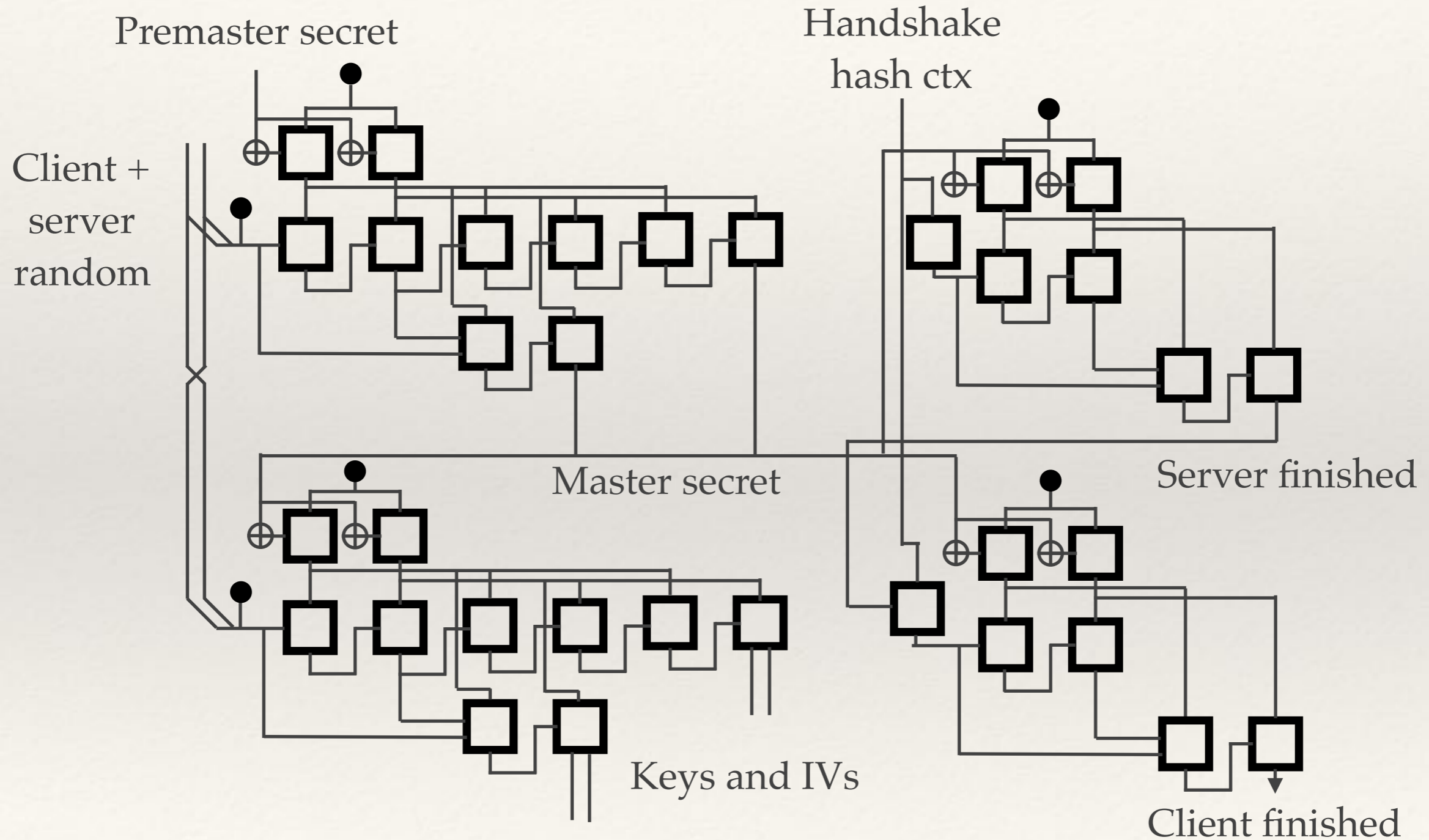
$$K_1 = \text{KDF}_1(\sigma, A, B, X, Y)$$

$$M_A = \text{MAC}(K_1; A, X)$$

$$M_B = \text{MAC}(K_1; B, Y)$$

$$K_2 = \text{KDF}_2(\sigma, A, B, X, Y)$$

Motivation: [TLS 1.2]



Finished is also encrypted, but I got bored before drawing the cipher calls.

The modern solution

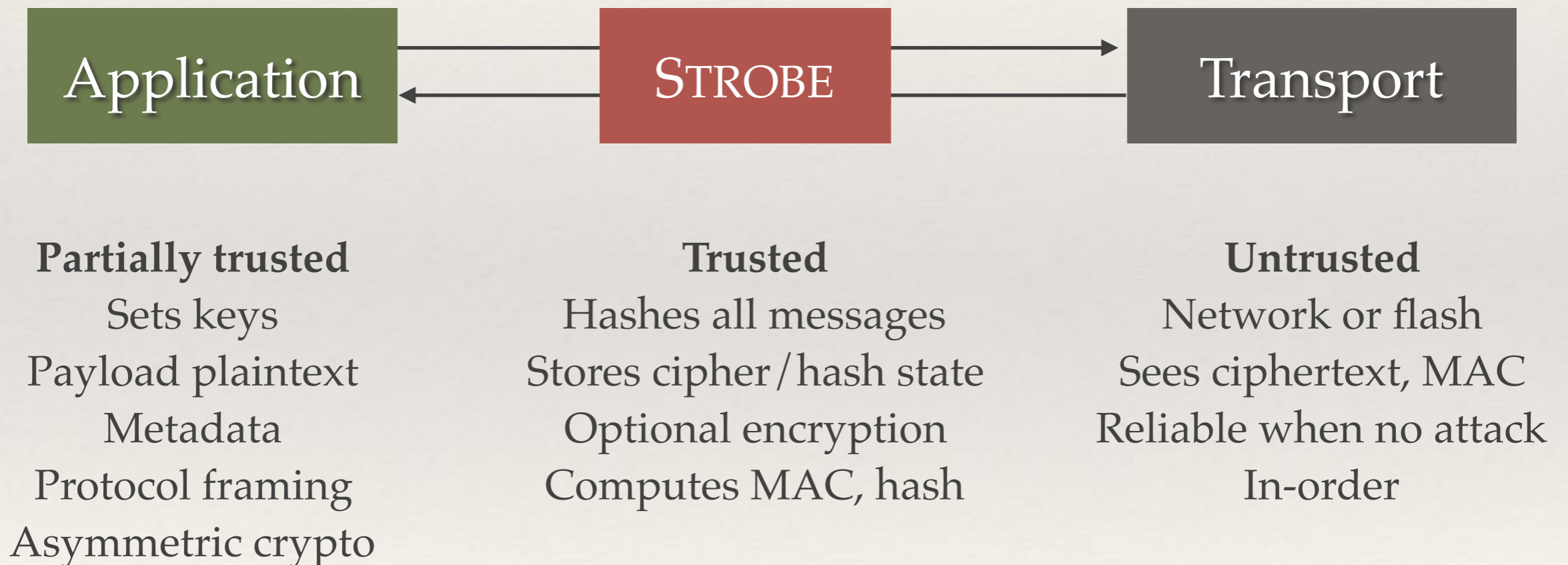
Hash all the things!



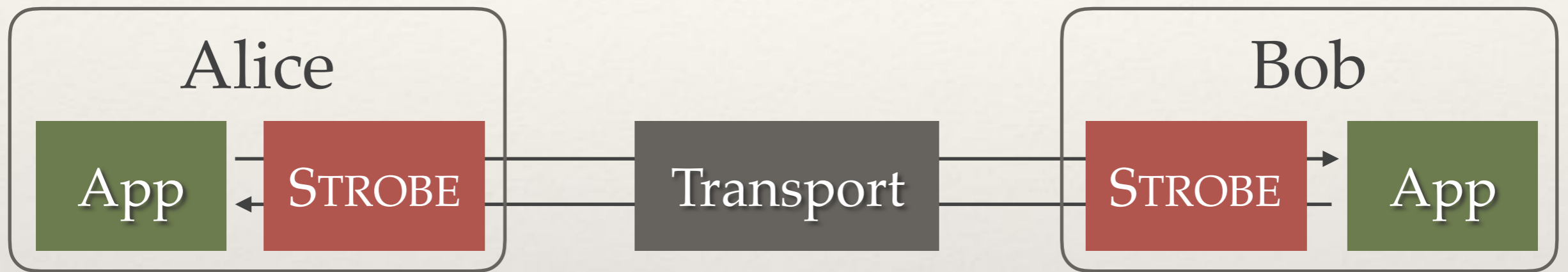
Eg: [TLS 1.3], [Noise], [BLINKER]

STROBE overview

All messages pass through STROBE
(at the least, to update the hash)



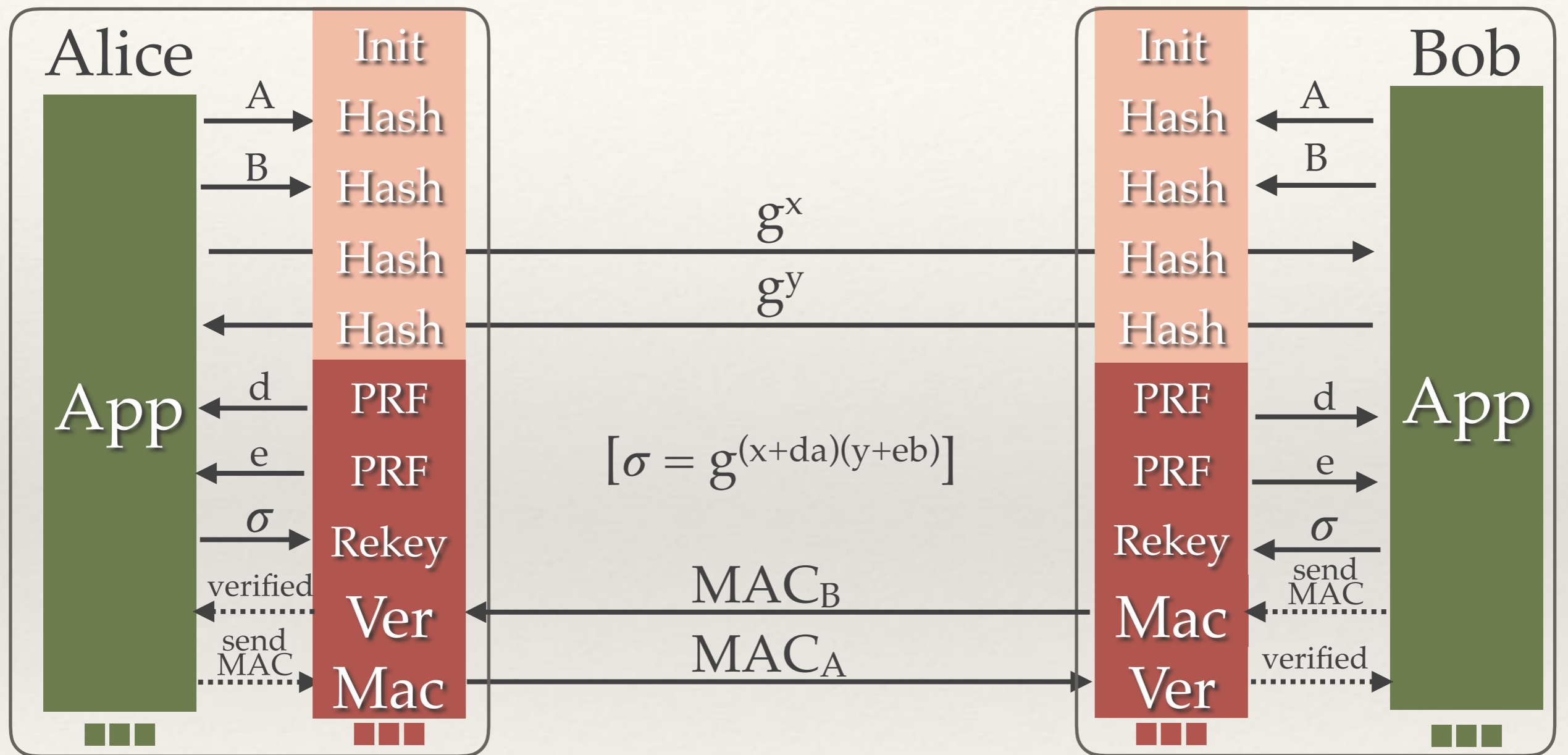
STROBE two-party protocols



Alice and Bob's STROBE instances advance in lockstep

If a message is changed on the wire, the next MAC will fail

STROBE example: FHMV-C



Everything is based on running hash

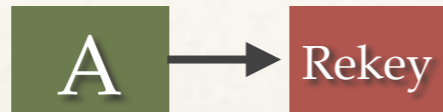
$$d = H(A, B, g^x, g^y) \quad e = H(A, B, g^x, g^y, d)$$

$$MAC_B = H(A, B, g^x, g^y, d, e, \sigma)$$

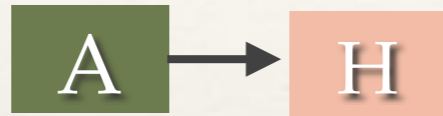
(roughly)

STROBE operations

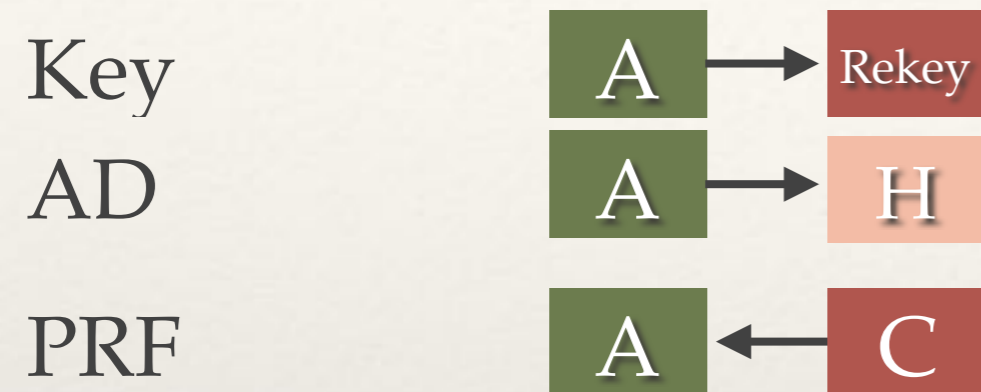
Key



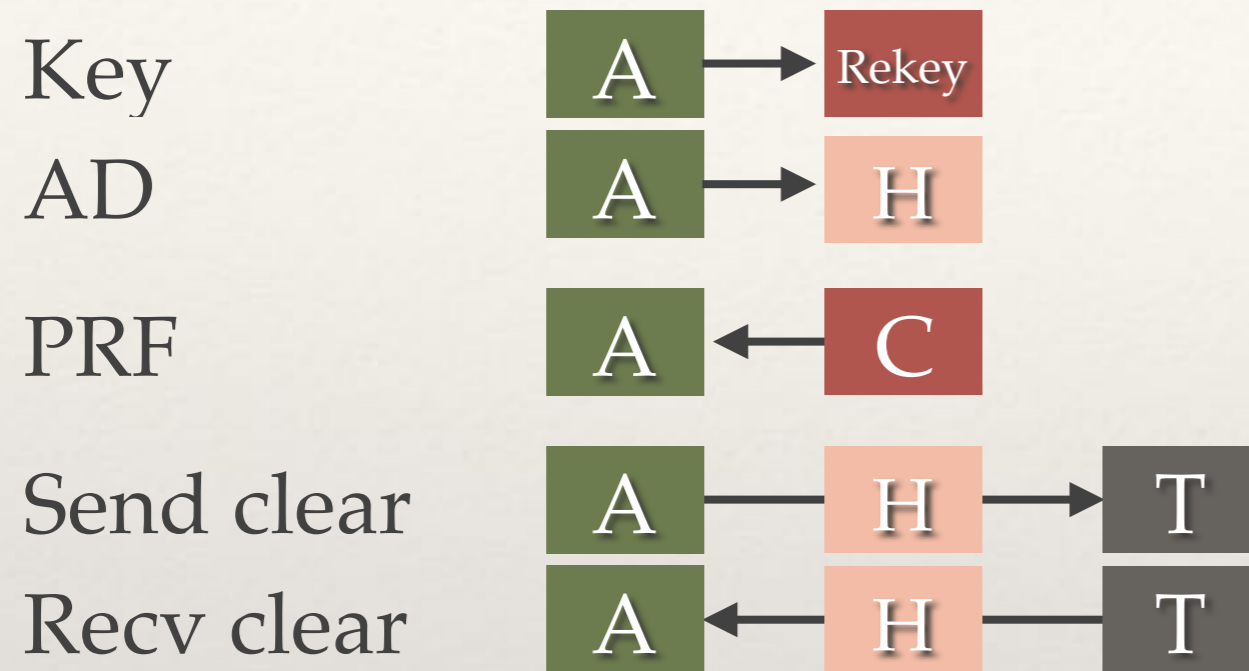
AD



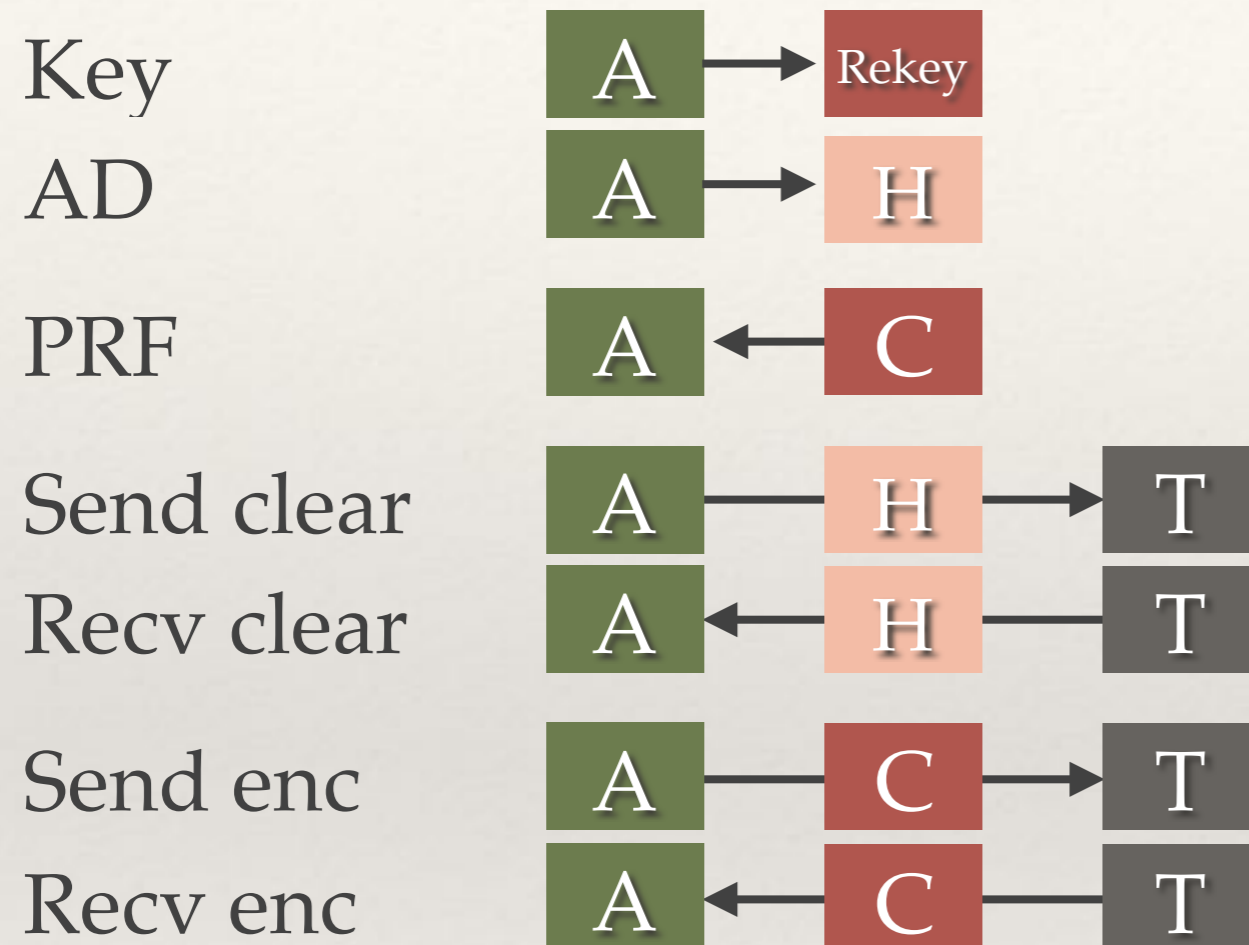
STROBE operations



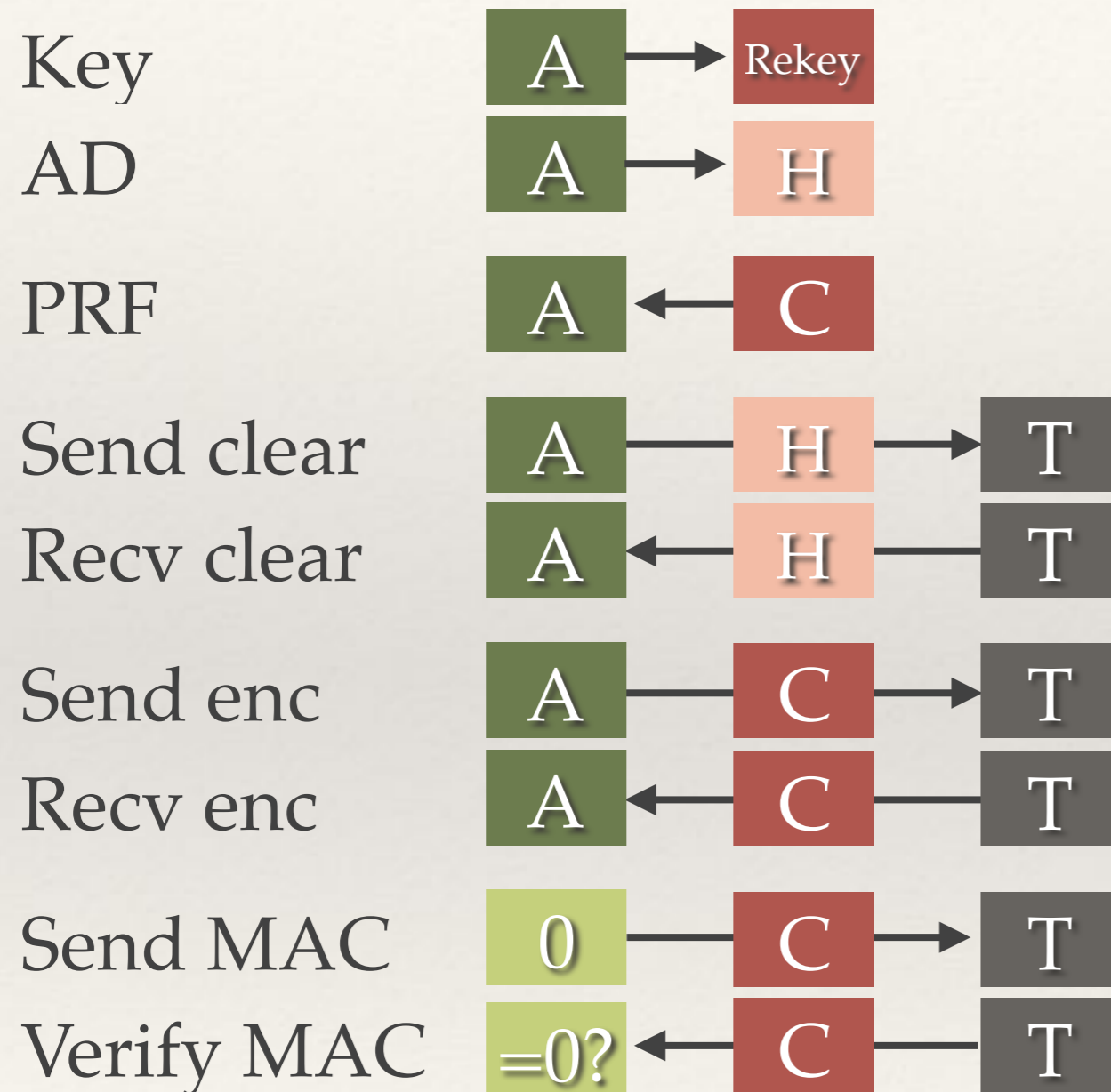
STROBE operations



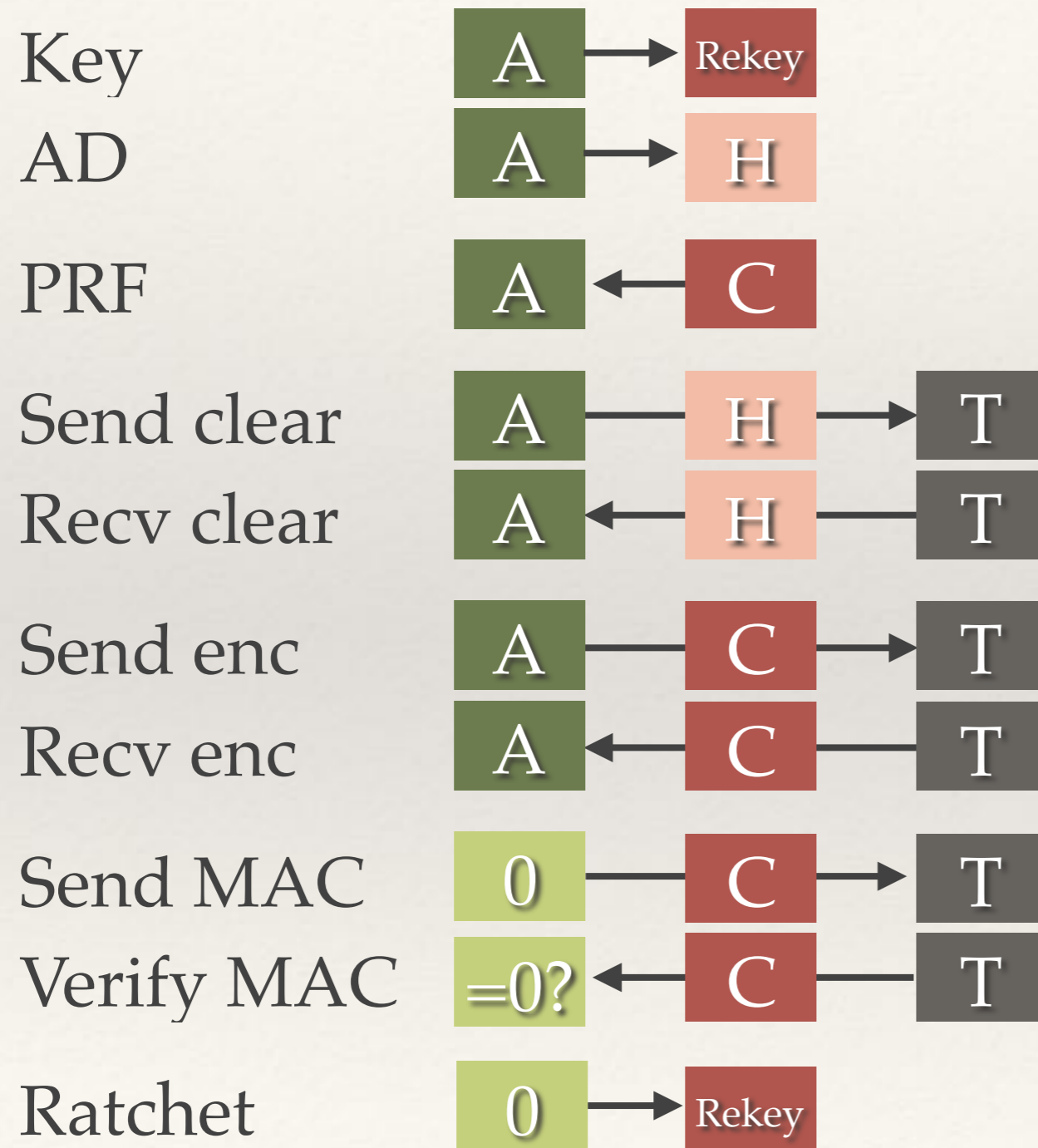
STROBE operations



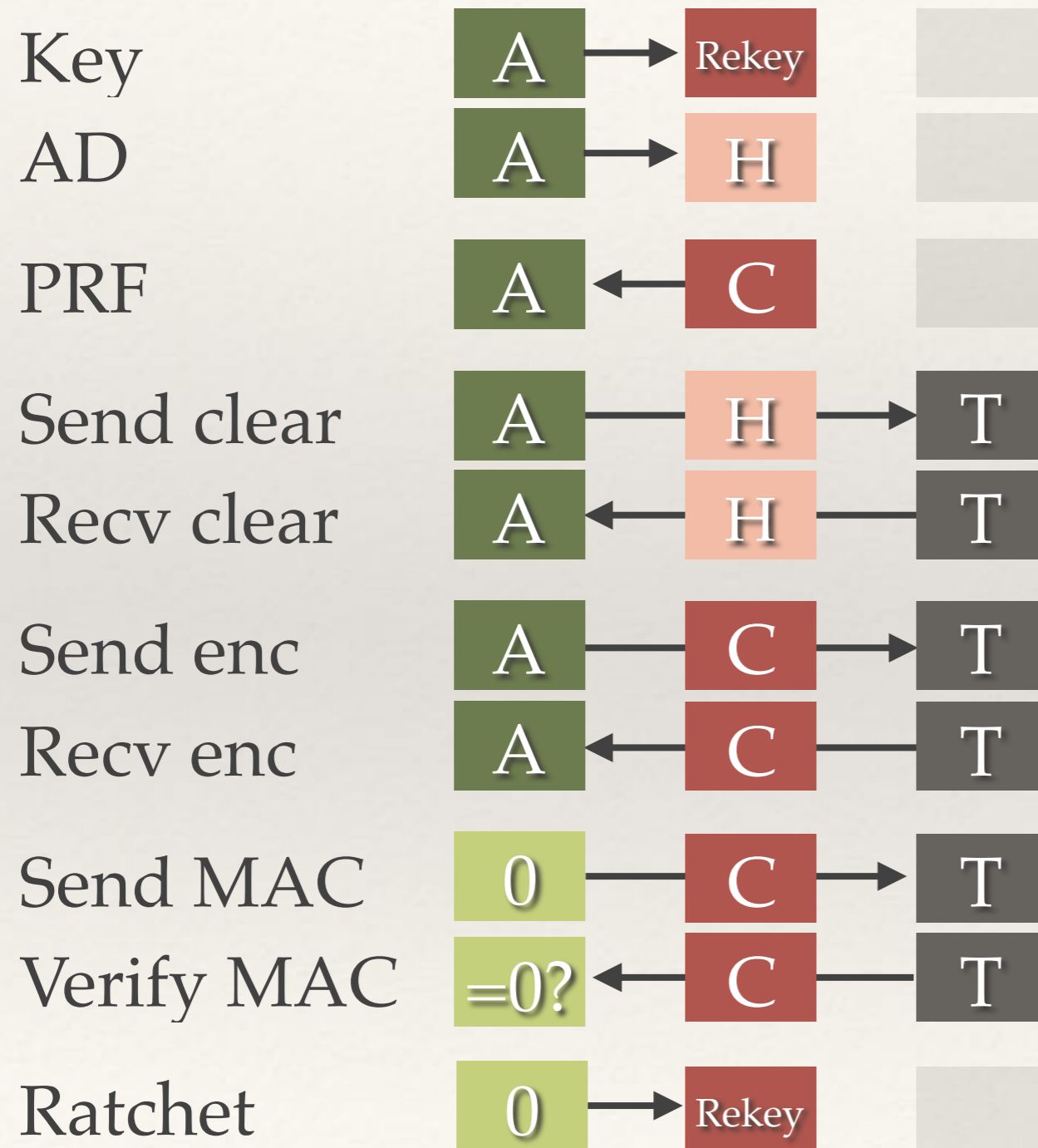
STROBE operations



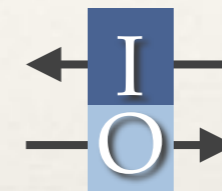
STROBE operations



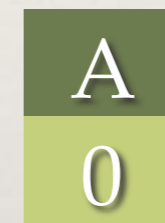
STROBE operations



All described by 4 features:



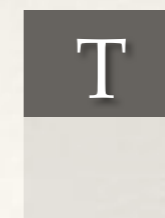
Data flow direction



Data goes to/
from app



Data goes to/
from cipher
(else just hash)

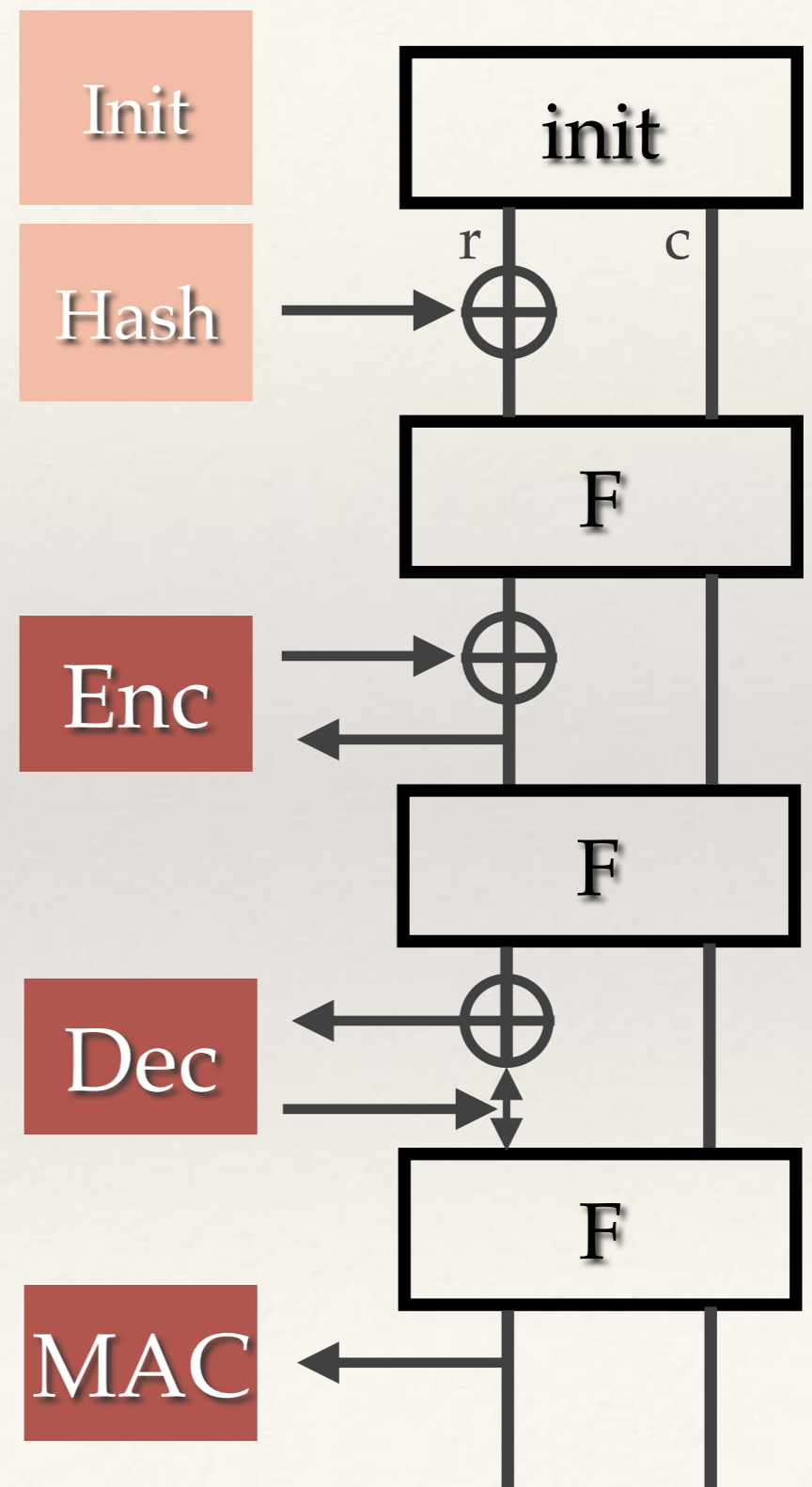


Data goes to/
from transport

STROBE implementation

- ❖ Duplex sponge construction
[RadioGatún, KECCAK, Duplex]

- ❖ State is divided into two parts:
 - ❖ Rate gets xor'd with input block
 - ❖ Capacity is kept separate
 - ❖ $(r, c) = F(r \oplus m, c)$



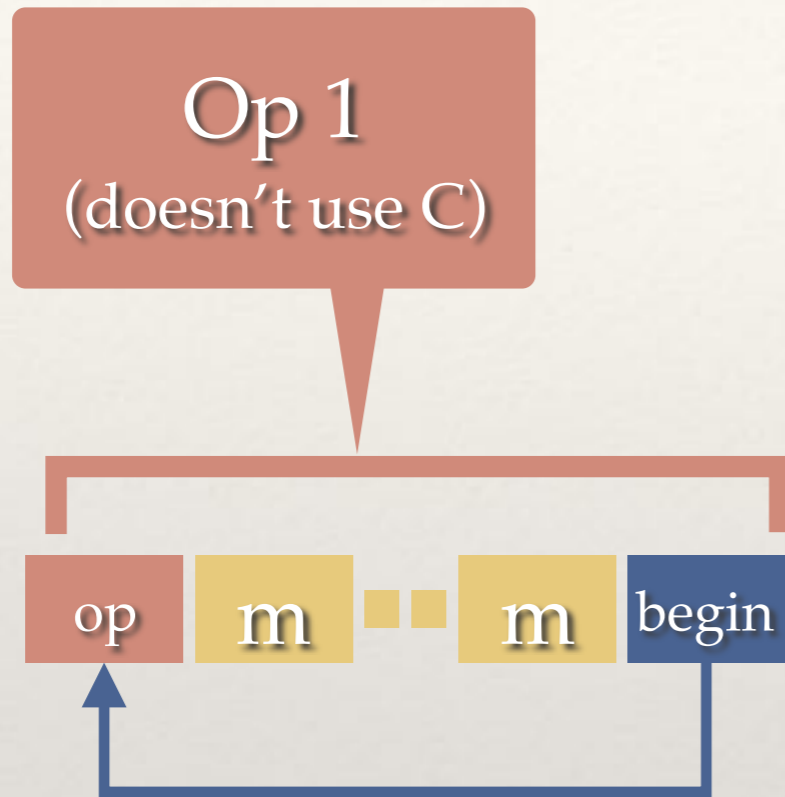
“Hash all the things”

- ❖ Goal: output of Strobe is a **random oracle**
- ❖ Input is all previous operations
 - ❖ $H(\text{“abc”}) \neq H(\text{“a”}, \text{“bc”})$
 - ❖ Includes operation type and data
 - ❖ Includes intended use of output

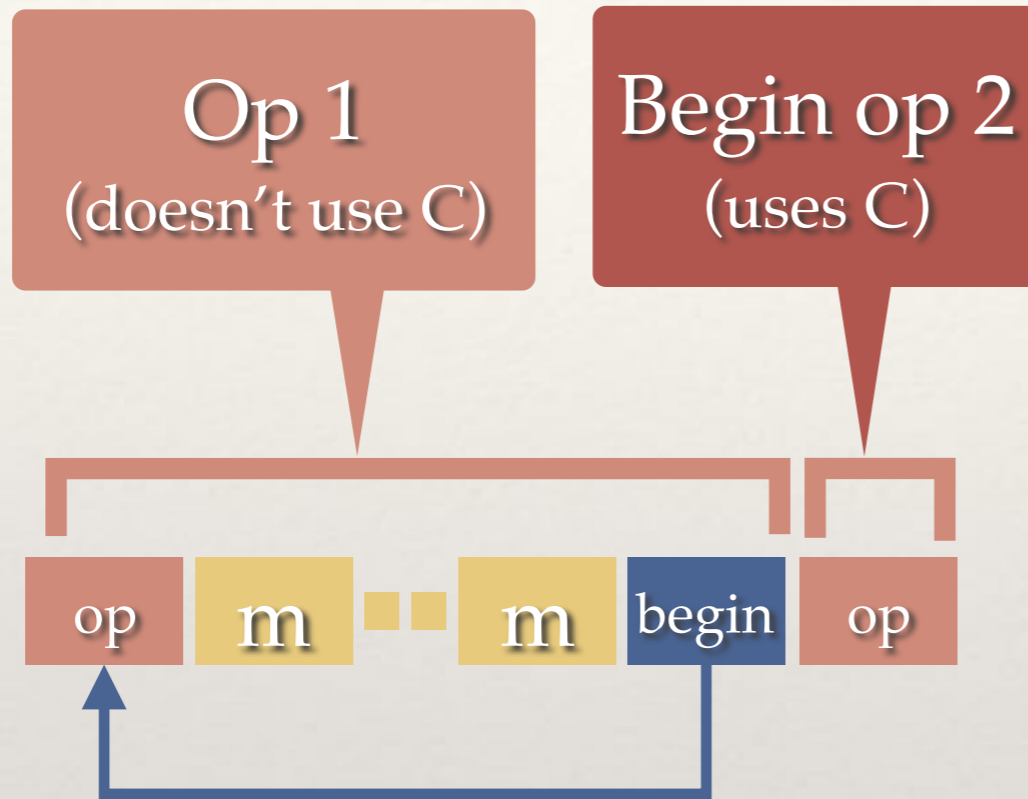
STROBE padding

- ❖ Theorem from [Duplex]: sponge output is a **random oracle** on previous inputs (if F is a random fn / perm)
- ❖ \Rightarrow Requirement: each time F is called, can parse:
 - ❖ Entire previous transcript
 - ❖ Intended use of output

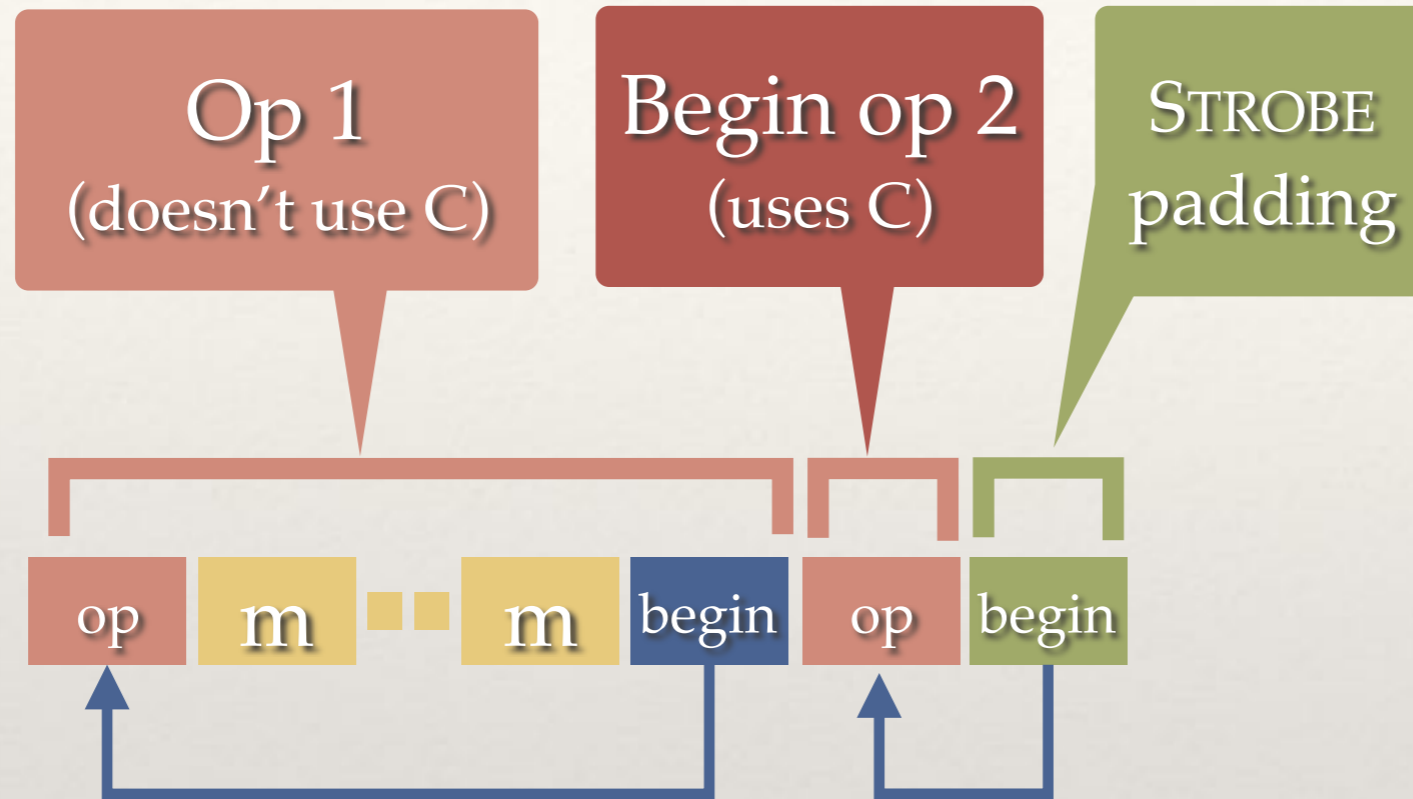
STROBE padding



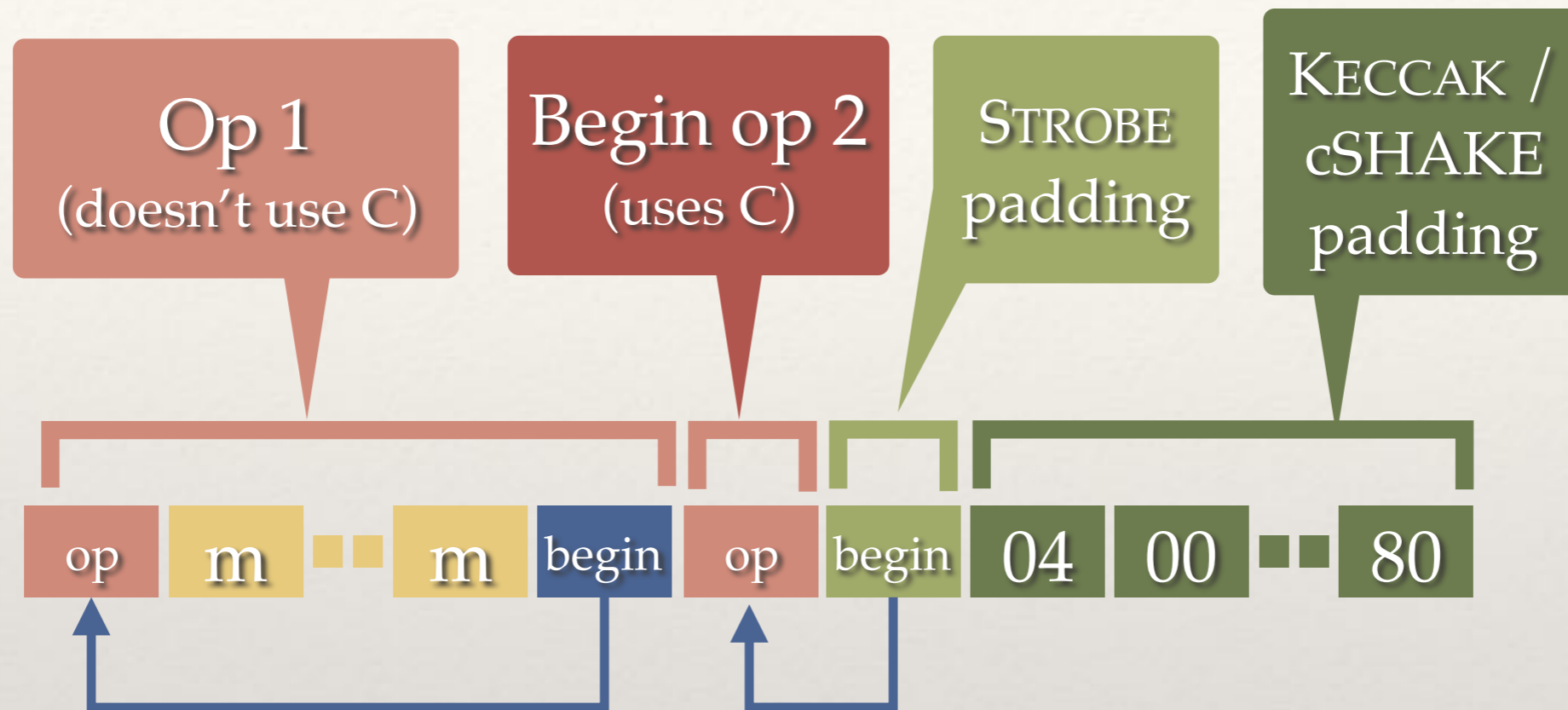
STROBE padding



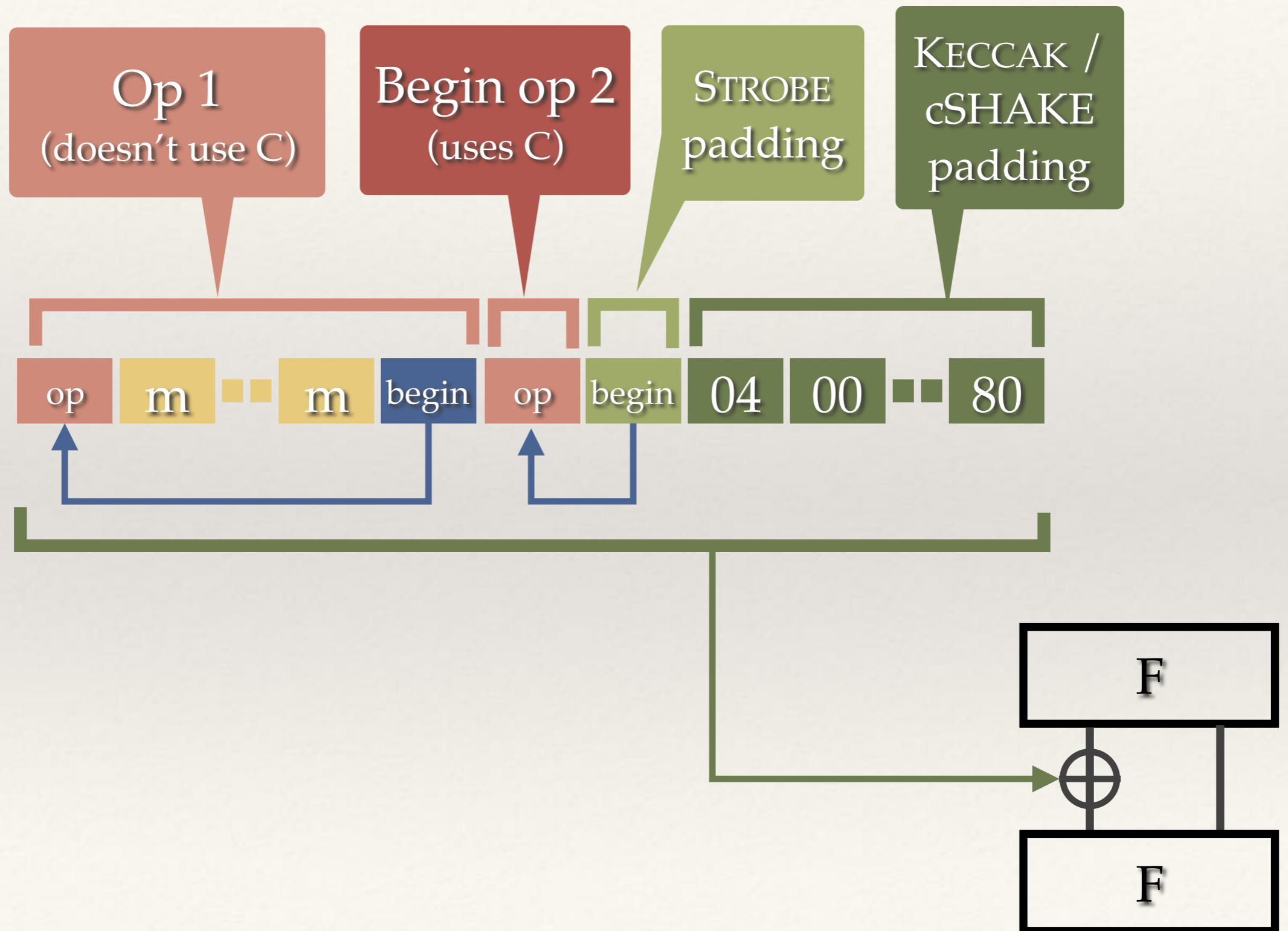
STROBE padding



STROBE padding



STROBE padding



Operations with metadata

- ❖ Output depends on its **intended use**
 - ❖ “Will be used to encrypt a message” isn’t good enough
 - ❖ What kind of message? How long?
- ❖ Disambiguate with metadata operations
 - ❖ Metadata AD / CLR / ENC before each operation
 - ❖ Can be (tag, length) of protocol framing
 - ❖ Optional but recommended. Cheap.

Implementation

- ❖ Prototype C code at <https://strobe.sourceforge.io/>
- ❖ Optimized for size on embedded devices
- ❖ Includes simple callback-based IO engine
- ❖ Curve25519 code may be of independent interest

Implementation results

- ❖ KECCAK- f [800], Cortex-M3 / M4 C
 - ❖ < 2 KB code, <350 B stack
- ❖ With Curve25519 ECDH / sign / verify; PRNG support
 - ❖ < 3.5KB code, 700B stack, 120B PRNG pool
- ❖ Significantly smaller with asm intrinsics (unreleased)

Future work

- ❖ Better documentation and example protocols
- ❖ Improve engine code
- ❖ Non-sponge implementation
- ❖ Formal analysis
 - ❖ Most work is done by [Duplex]
 - ❖ Rollback resistance, full protocol analysis
- ❖ Post-quantum analysis

Works cited

- ❖ [BLINKER]: Markku-Juhani Saarinen. “Beyond Modes: Building a Secure Record Protocol from a Cryptographic Sponge Permutation.” CT-RSA 2014, <https://eprint.iacr.org/2013/772>
- ❖ [cSHAKE]: John Kelsey, Shu-jen Chang Ray Perlner. “SHA-3 Derived Functions: cSHAKE, KMAC, TupleHash and ParallelHash.” NIST SP 800-185, December 2016, <http://nvlpubs.nist.gov/nistpubs/SpecialPublications/NIST.SP.800-185.pdf>
- ❖ [Duplex]: Guido Bertoni, Joan Daemen, Michaël Peeters, Gilles Van Assche. “Duplexing the sponge: single-pass authenticated encryption and other applications.” SAC 2011, <http://sponge.noekeon.org/SpongeDuplex.pdf>
- ❖ [FHMQV]: Augustin Sarr, Philippe Elbaz-Vincent, Jean-Claude Bajard. “A secure and efficient authenticated diffie–hellman protocol.” European PKI Workshop 2009, <https://eprint.iacr.org/2009/408>

Works cited

- ❖ [KECCAK]: Guido Bertoni, Joan Daemen, Michaël Peeters, Gilles Van Assche. “The KECCAK sponge function family.” NIST SHA-3 submission, <http://keccak.noekeon.org/>
- ❖ [Noise]: Trevor Perrin. “Noise Protocol Framework.” <http://www.noiseprotocol.org/>
- ❖ [RadioGatún]: Guido Bertoni and Joan Daemen, Michaël Peeters and Gilles Van Assche. “RadioGatún, a belt-and-mill hash function.” Cryptographic Hash Workshop 2006, <http://eprint.iacr.org/2006/369>
- ❖ [TLS 1.2]: Tim Dierks and Eric Rescorla. “The Transport Layer Security (TLS) Protocol, Version 1.2.” RFC 5246 (2008), <https://www.ietf.org/rfc/rfc5246.txt>
- ❖ [TLS 1.3]: Eric Rescorla. “The Transport Layer Security (TLS) Protocol Version 1.3.” draft, <https://tlsWG.github.io/tls13-spec/>

FIN

Questions?